Post-quantum Resettably-Sound Zero Knowledge

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Post-quantum Resettably-Sound ZK

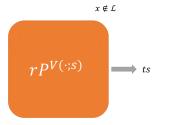
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Quantum Resettable Soundness

We initiate the study of Quantum Resettable Soundness

• Attackers perform *quantum* resetting attacks.



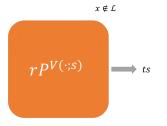
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• Attackers perform *quantum* resetting attacks.



- Motivations:
 - Achieving such a security guarantee.
 - Better understand post-quantum black-box zero knowledge.

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Lower-Bounds on Post-quantum Black-Box Zero Knowledge

- Post-quantum zero knowledge protocols that are **three message or constant-round public-coin** cannot be black-box zero knowledge, except for languages in **BQP**.
 - Classical lower-bound don't extend to the quantum case simulator has more power.

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Observation

If a language \mathcal{L} has a post-quantum black-box zero-knowledge, resettably sound protocol, then $\mathcal{L} \in BQP$.

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Observation

If a language \mathcal{L} has a post-quantum black-box zero-knowledge, resettably sound protocol, then $\mathcal{L} \in BQP$.

- Methodology Transform protocol into resettably sound one.
 - Using measure-and-reprogram techniques (Don-Fehr-Majenz(20)[5]).

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Constructing a Quantum Resettably Sound Zero Knowledge Protocol

We present a construction of a *post-quantum* resettably sound zero knowledge protocol for **NP**.

- Assumes QLWE and quantum FHE.
- The construction is constant round.
- Zero knowledge holds with regards to quantum verifiers.
- Starting point The protocol of Bitansky-Shmueli(20)[4].
 - As is, protocol is not resettably sound, even classically.

From Resettable Soundness to Quantum Unobfuscatable Functions

A connection between resettable soundness and quantum VBB obfuscation.

- We show that if there exists a post-quantum resettably-sound zero-knowledge argument for **NP** and post-quantum one-way functions, then quantum VBB obfuscation is impossible.
- An analog of the classical result of Bitansky-Paneth(15)[3].

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- An analog of the classical result of Bitansky-Paneth(15)[3].
- An alternative proof to results from Alagic-Brakerski-Dulek-Schaffner(20)[1], Ananth-La Placa(20)[2].
- Better resettably sound protocols imply better construction of unobfuscatability.

- G. Alagic, Z. Brakerski, Y. Dulek, and C. Schaffner. Impossibility of quantum virtual black-box obfuscation of classical circuits. *CoRR*, abs/2005.06432, 2020.
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